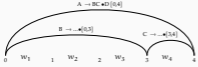


- We can formulate parsing as
  - Top-down: begin with the start symbol, try to produce the input string to be parsed
  - Bottom-up: begin with the input, and try to reduce it to the start symbol
- For both options, we have seen examples of chart parser
- Parsing can also be *directional* or *non-directional*
- In this lecture, we introduce a general mechanism for chart parsing that has all these forms of parsing methods as special cases

The overall idea

- We adopt Early-like chart entries of the form:  $X \rightarrow \alpha \cdot \beta$  [ $i, j$ ] where,
  - $i$  and  $j$  are indexes starting from 0 (0 indicating beginning of the input string)
  - The chart entry indicates  $\alpha$  is found between  $i$  and  $j$ , we are looking for a  $\beta$  starting from  $j$
- At any time, we have two sets of items:
  - active items are those we expect to complete
  - inactive items are those with a dot at the end
- The goal is to complete  $S \rightarrow \cdot \alpha$  [ $0, n$ ]



Components of a typical chart parsing algorithm

- Besides the chart, we keep an agenda of 'unexplored items'
- A set of inference rules determine how to modify the chart when processing items from the agenda
- Typically inference rules are similar to completion process of Earley parser
- The following inference rule is part of every chart parser (so-called 'fundamental rule' of chart parsing)
  - If there is an inactive item of the form  $A \rightarrow \alpha \cdot$  and an active item of the form  $B \rightarrow \beta \cdot \alpha y$  add item  $B \rightarrow \beta \cdot y$
- We also need a strategy for selecting the items from the agenda and applying the inference rules
- Depending on the data structure used for the agenda, and order of processing of inference rules, we may get different types of parsers

The sketch of a chart parsing algorithm

```

1: Initialize A (agenda) and C (chart)
2: repeat
3:   i ← next(A)
4:   if i ∈ C then
5:     discard i
6:   else
7:     apply all inference rules to i
8:     place new items in A
9:     place the item in C
10: until A is empty
    
```

- Very simple, but unspecified parts:
  - Initialization
  - Inference rules
  - The order of items received from the agenda
- An item is put into chart only after all inferences from it are in the chart or in the agenda
- Chart is a set, items do not repeat

Bottom-up chart parsing

- Single additional inference rule:
  - If a new item has the form  $A \rightarrow \alpha \cdot$ , add  $B \rightarrow A \cdot \beta$  for each rule  $B \rightarrow A \beta$  in the grammar.
- Initialization:
  - Empty chart
  - Place  $P \rightarrow w_i$  [ $i-1, i$ ] in the agenda for all word  $w_i$  ( $P$  is the pre-terminal symbol, typically the PDS tag in CL)
  - If there are  $\alpha$  rules, add  $P \rightarrow \cdot \alpha$  [ $1, i$ ] for all  $P \rightarrow \alpha$  in the grammar, for  $i \in \{0, n\}$
- Choice of agenda does not matter. A stack is typical, but a queue or a priority queue is also an option

Example: bottom-up chart parsing

grammar

```

S → NP VP
NP → Pn N
NP → Ptn
VP → V NP
VP → V
Ptn → I
Ptn → her
V → saw
N → duck
N → .
    
```

stack

Ptn → I [0,1]
V → saw [1,2]
Ptn → her [2,3]
N → duck [3,4]

Example: bottom-up chart parsing

grammar

```

S → NP VP
NP → Pn N
NP → Ptn
VP → V NP
VP → V
Ptn → I
Ptn → her
V → saw
N → duck
N → .
    
```

stack

NP → Ptn N [0,1]
NP → Ptn N [0,1]
V → saw [1,2]
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NP → Pn N
NP → Ptn
VP → V NP
VP → V
Ptn → I
Ptn → her
V → saw
N → duck
N → .
    
```

stack

I → NP VP [0,1]
NP → Ptn N [0,1]
V → saw [1,2]
Ptn → her [2,3]
N → duck [3,4]

Example: bottom-up chart parsing

grammar

```

S → NP VP
NP → Pn N
NP → Ptn
VP → V NP
VP → V
Ptn → I
Ptn → her
V → saw
N → duck
N → .
    
```

stack

NP → Ptn N [0,1]
V → saw [1,2]
Ptn → her [2,3]
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Example: bottom-up chart parsing

grammar

```

S → NP VP
NP → Pn N
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Ptn → I
Ptn → her
V → saw
N → duck
N → .
    
```

stack

V → saw [1,2]
Ptn → her [2,3]
N → duck [3,4]

Example: bottom-up chart parsing

grammar

```

S → NP VP
NP → Pn N
NP → Ptn
VP → V NP
VP → V
Ptn → I
Ptn → her
V → saw
N → duck
N → .
    
```

stack

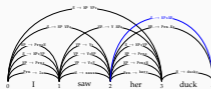
NP → Ptn N [1,2]
VP → V NP [1,2]
VP → V [1,2]
Ptn → her [2,3]
N → duck [3,4]



## Example: bottom-up chart parsing

grammar

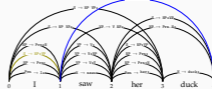
$S \rightarrow NP VP$   
 $NP \rightarrow Prn N$   
 $NP \rightarrow Prn VP$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V$   
 $VP \rightarrow V S$   
 $Prn \rightarrow I$   
 $Prn \rightarrow her$   
 $V \rightarrow saw$   
 $N \rightarrow duck$   
 $V \rightarrow duck$



## Example: bottom-up chart parsing

grammar

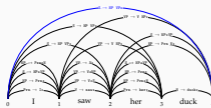
$S \rightarrow NP VP$   
 $NP \rightarrow Prn N$   
 $NP \rightarrow Prn VP$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V$   
 $VP \rightarrow V S$   
 $Prn \rightarrow I$   
 $Prn \rightarrow her$   
 $V \rightarrow saw$   
 $N \rightarrow duck$   
 $V \rightarrow duck$



## Example: bottom-up chart parsing

grammar

$S \rightarrow NP VP$   
 $NP \rightarrow Prn N$   
 $NP \rightarrow Prn VP$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V$   
 $VP \rightarrow V S$   
 $Prn \rightarrow I$   
 $Prn \rightarrow her$   
 $V \rightarrow saw$   
 $N \rightarrow duck$   
 $V \rightarrow duck$



## Bottom-up chart parsing

## additional remarks

- The parser (as described) proceeds bottom up (left-corner)
- It can process arbitrary CF grammars
- Stack-based agenda is common, queue-based agenda is rarely used
- An interesting alternative is so-called *head-corner* parsing: using a priority queue (e.g., processing 'heads' first) one can build the 'important' parts of the tree first
- The time complexity is  $O(n^3)$
- There are many variants, optimizations (based on, different inference rules, processing strategies)

## Top-down chart parsing

- The basic algorithm is the same, but we specify
  - Inference rule (besides the 'fundamental rule'):
    - If the new edge has the form  $A \rightarrow \alpha \beta B$  [ $s, j$ ], add  $B \rightarrow \gamma$  [ $j, j$ ] for each rule  $B \rightarrow \gamma$  in the grammar.
- Initialization
  - Empty chart
  - Push  $\rightarrow \epsilon$  [ $0, 0$ ] into the stack
  - Push all productions for the terminal symbols to the stack (or to the chart, as there is nothing to predict for these productions)
- Typically we use a stack as an agenda

## Example: top-down chart parsing

grammar

$S \rightarrow NP VP$   
 $NP \rightarrow Prn N$   
 $NP \rightarrow Prn VP$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V$   
 $VP \rightarrow V S$   
 $Prn \rightarrow I$   
 $Prn \rightarrow her$   
 $V \rightarrow saw$   
 $N \rightarrow duck$   
 $V \rightarrow duck$



## Top-down chart parsing

## additional remarks

- The parser (as described) is purely top-down
- In practice, it is common to use 'lookup'
- Stack-based agenda is common, queue-based agenda is rarely used
- The time complexity is  $O(n^3)$
- There are many variants, optimizations (based on, different inference rules, processing strategies)

## Summary

- Chart parsing is a general framework for constructing a variety of parsers
- It shares many similarities with the CKY and Earley

## Next:

- Deterministic parsing (maybe after the break)

## Acknowledgments, references, additional reading material

[https://doi.org/10.1007/978-1-4939-9826-2\\_11](https://doi.org/10.1007/978-1-4939-9826-2_11) [http://dx.doi.org/10.1007/978-1-4939-9826-2\\_11](http://dx.doi.org/10.1007/978-1-4939-9826-2_11)

### Example: top-down chart parsing

```
grammar
S → NP VP
NP → Pm N
VP → V NP
VP → V
Pm → I
Pm → her
V → saw
N → duck
V → duck
```

0 I 1 saw 2 her 3 duck 4

